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Title: **DATA PRIVACY PRESERVATION AND LOCATION OF A ALTERNATE USER IN A WIRELESS COGNITIVE NETWORKS**

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Paper Authors

M.BHADRAJA, E.V.V.S.SIVAKUMAR , S.RAJESH

PB Siddhartha College of arts and science, vijayawada



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DATA PRIVACY PRESERVATION AND LOCATION OF A ALTERNATE USER IN A WIRELESS COGNITIVE NETWORKS

¹M.BHADRAJA, ²E.V.V.S.SIVAKUMAR, ³S.RAJESH

¹Lecturer, Department of Computer Science, PB Siddhartha College of arts and science, Vijayawada.

²Lecturer, Department of Computer Science, PB Siddhartha College of arts and science, Vijayawada.

Lecturer, Department of Computer Science, PB Siddhartha College of arts and science, Vijayawada.

ABSTRACT: The extant system provides location privacy preserving schemes for cognitive radio networks (CRNs) that protect secondary users (SUs) location privacy while allowing them to preserve their own sensitive details. It contains harness probabilistic set membership data structures to exploit the structured nature of spectrum databases (DBs) and SUs queries. This enables us to create a compact representation of DB that could be queried by SUs without having to share their location with DB, thus guaranteeing their location and data privacy. In our proposed work, the system allows user to register in the Network and allows them to upload any kind of sensitive detail through setting down into it. The system provides security for the data of a secondary user (SU) in a cognitive network through AES encryption algorithm. If any of the intruder in network tries to steal user's details, they will use fake details for setting down. In the mean time, server in the network is warned about cybercriminal entering and it tracks IP address, ISP and geographical information of the intruder and blocks invader from entering within the network.

KEYWORDS: Privacy preservation, Wireless Cognitive Network, Advanced Encryption Standard (Algorithm), LPDBQS (Algorithm), Interloper attack

1. INTRODUCTION

A **wireless network** is a computer network that uses wireless data connections between network nodes. Wireless networking is a method by which homes, telecommunications networks and business installations avoid the costly process of introducing cables into a building, or as a connection between various equipment locations (Fig.1).

Wireless network classified as:

Wireless PAN (personal area network)

Wireless LAN (local area network)

Wireless ad hoc network

wireless MAN (metropolitan area networks)

Wireless WAN (wide area networks)

Cellular network

Global area

network Space

network



Fig. 1. Wireless Network

Wireless Cognitive Network(WCN):

Cognitive radio is an adaptive, intelligent radio and network technology that can automatically detect available nodes in a wireless spectrum and change transmission parameters enabling more communication to run concurrently. Cognitive Radio Network (CRN) is regarded as an emerging technology to address the increasing demand for node resources. It solves the node resource shortage problem by allowing a Secondary User (SU) to access the channel of a Primary User (PU) when the channel is not occupied by the PU, in which an SU queries a database to obtain node availability information by submitting a location based query (Fig.2).

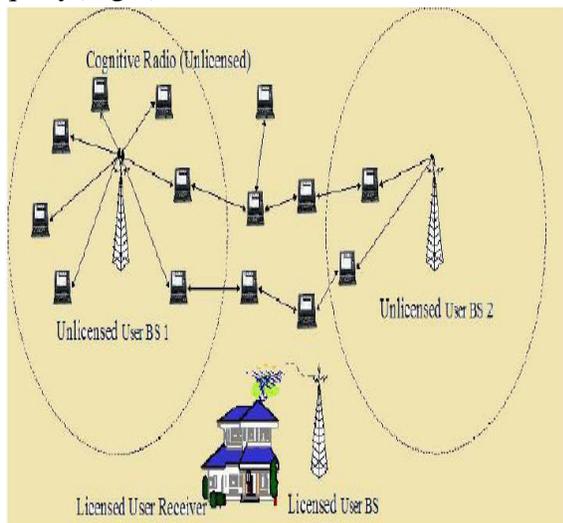


Fig.2. Wireless Cognition Network

Primary user (PU): They were the **licensed user** of the network and have the access for their account features.

Secondary user (SU): They were the **unlicensed user** of the network. Also known as Mobile users (on travel).

In CRNS, the SUs are allocated with primary user's node for accessing when they are at rest from network. In this case, the data of the SU were visible to PU. For this issue

many researches were undertaken and found solutions by introducing two party protocols, filters such as cuckoo and bloom. Though the security is provided for SUs the intruder's in the network tries to hack the network and steals the data of the users. Security is a major factor that reduces the performance in Wireless Cognitive Network (WCN). However, one concern about database-driven CRNs is that the queries sent by SUs will inevitably leak the location information. Instead of directly learning the SUs' locations from their queries, our discovered attacks can infer an SU's location through his used nodes. The location privacy preservation schemes for database driven cognitive radio networks provides an optimal location privacy to secondary users within database coverage network by leveraging set membership data structure to construct a compact version of database. Even though the location is preserved from intruder, they try to hack the sensitive details of SUs such as Business information, Personal information, Financial information etc.,.. This system allows the SUs to store the sensitive details and keep track from hackers. If any hacker tries to hack the network where the SUs work, the server in the network tracks the intruder and finds their IP address, ISP, Gateway and geological information of the SUs like latitude and longitude value. Finally blocks the Hackers.

Disadvantages of extant system:

- The location privacy issue in database-driven CRNs.
- It introduces some noise to SU's location which may cause erroneous spectrum availability information.

II. LITERATURE REVIEW

The authors Mohamed Grissa et al., in 2017 introduced a system that allows to preserve the location privacy of SUs while performing reliable and efficient spectrum sensing using Cryptographic mechanisms[4]. The researchers S. Selvakamani et al., in 2017 presented a system that achieves a balance by minimizing interference to licensed users and maximizing the entire system performance providing opportunistic access to number of secondary users such as opportunistic spectrum sensing and adaptive channel assignment through mathematical analysis and MC-OSACA technique [9]. The experts H. Zhu et al., in 2016 invented Jammer Inference based Jamming Defense (jDefender) Framework. The main idea of jDefender is inferring the likelihood of a user being a jammer based on the observed jamming events and then utilizing the inferred attack likelihood to enhance the effectiveness of a series of proposed anti-jamming strategies[2]. The researchers Mohamed Grissa, et al., in 2015 proposed a system that provides an efficient scheme for database driven CRNs that preserves the location privacy of SU through Cuckoo filter [1]. The experts Xu Zhang, et al., in 2014 provided a comprehensive analysis and guide of existing efforts around localization and location privacy preservation in cognitive radio network. The cognoscenti Z. Gao, et al., in 2013 designed Private Spectrum Availability Information Retrieval scheme that utilizes a blind factor to hide the location of the SU and proposed a novel prediction based Private Channel Utilization

protocol that reduces the possibilities of location privacy leaking by choosing the most stable channels[11]. Despite its importance, the location privacy issue in CRNs only recently gained interest from the research community[4]. Some works focused on addressing this issue in the context of collaborative spectrum sensing [5]–[8],[14]–[15]. Protecting SU's location privacy in database-driven CRNs is a very challenging task, since SUs are required to provide their physical locations to DB in order for them to be able to learn about spectrum opportunities in their vicinities[1]. However, direct adaptation of such concepts yield either insecure or extremely costly results. For instance, k -anonymity guarantees that SU's location is indistinguishable among a set of k points, which could be achieved through the use of dummy locations by generating k properly selected dummy points, and performing k location privacy and maximizing some utility, which makes it suffer from the fact that achieving a high location privacy level results in a decrease in spectrum utility. PIR, on the other hand, allows a client to obtain information from a database while preventing the database from learning which data is being retrieved. Several approaches have used this approach[11] proposed a PIR-based approach.

III. PROPOSED SYSTEM

A. System model:

The Cognitive radio network that consists of a set of Secondary users stores geo-location information on the database (DB). SUs are assumed to be enabled with GPS and node sensing capabilities, and to have access to DB to obtain node availability information

within its operation locale. The main issue faced by the SUs in CRNs is the state of being free and cybercriminal attacks. The proposed system solves this issues by storing the details of along with location in the database using Encryption algorithms. When the Intruder tries use the SUs details, the presented system immediately notices the IP address, ISP of the Intruder, Gateway and geological informations such as latitude and longitude values. When the server in the CRN knows about the attack as soon as he blocks them. So that the Intruder won't be able to get into that network through that system.

B. System Design:

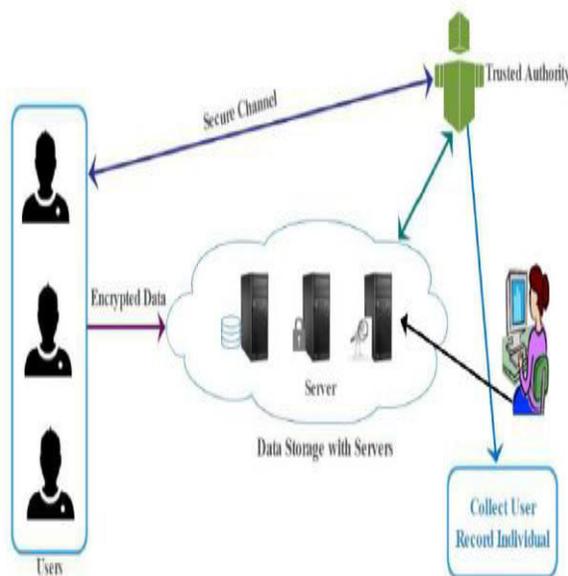


Fig.3. Overview of extant system

This fig:1. Describes how the authority member provides channels for the SUs and they were using the network by querying the network. The authority member stores the information of the user in a separate database that is visually hidden from the user

but abstracted. And also it shows how the Intruder tries to hack the network details.

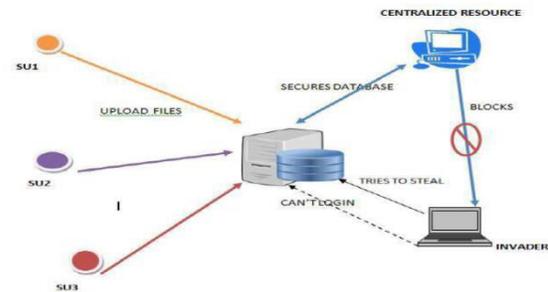


Fig.4: Proposed System

C. Detailed Description:

- Network ID Creation and Uploading data.
- Examine the Intruder and their IP, ISP and Geographical information.
- Blocking of User

Network ID creation and Uploading Data

In this , the users in the network registers their own details like user name and password for Network ID registration . This registration will track up location of the secondary user and automatically acquires geo-graphical informations such as latitude and longitude values and send those back to the Centralized resource. The server of the network allows users to upload file and stores their details in a database. While uploading data of the user in network , the system generates a key for file using Encryption algorithms. If user wants those details for their use, they will be allowed to download files from web at anytime using the key .

Examine the intruder and their ip, isp and geographical information

If hacker tries to access the network illegally, then they will uses fake user name

and password to get in the network. It will show incorrect user name details. In the mean time, the server is warned about malicious entering inside the network. Then they check details of the interloper. Once they get the details of invader, the centralized person tracks details of invader such as Internet Protocol address, Internet service provider's, Gateway and Geographical information such as latitude and longitude. By referring geographical information, the server gets their location and also able to view the invader's Location through maps.

Blocking Of User

In this part of the proposed system, When centralized person realized about malicious activity in network by invader, they identifies the geo-location and network providers of them. Immediately they will block invader system's IP address from network, so that this people will not be able to set in to the network using their system in anyway.

D. Algorithm Used:

1. LPDBQS (Location Privacy DataBase Query Server)

```

1: SU queries DB with query f(k; char; ts);
2: DB retrieves resp containing r entries satisfying char;

3: DB constructs CFk;
4: for j = 1; : : : r do
5: if avl j = 1 then
6: x j (locX jklocY jktsk : : : krowj(c));
7: CFk: InsertHMACk (x j);
8: DB sends CFk to QS over a high throughput link;
9: SU initializes decision Channel is busy
10: for all possible combinations of par do

11: SU computes y (locX klocY kchnktsk : : : kparn);
12: SU computes yk HMACk(y) and sends it to QS;
13: QS looks up for yk in CFk using Lookup;
14: if CFk:Lookup(yk) then
15: SU senses chn;
16: if Sensing(chn) available then

```

17: decision chn is available; break return decision LPDBQS does not leak any information about SU s' location beyond HMAC secure values. LPDBQS, which offers better performance at SU s' side than

that of existing system algorithm. Here the proposed offers better performance at SU s' side than that of LPDB in extant system. This comes at the cost of deploying an additional entity, referred to as query server (QS), and having a computational security as opposed to unconditional. QS is introduced to handle SU s' queries instead of DB itself, which prevents DB from learning information related to SU s' location information. QS learns nothing but secure messages sent by SUs to check the availability of a specific channel.

STEPS IN LPDBQS:

STEP 1

SU queries the database about the availability nodes. The database contents is retrieved as a collection of CF (only the entries that have available channels) by sending a secret key k.

STEP 2

DB sends CFk to QS over a high throughput link

STEP 3

SU read about the channel engaging and tries to adapt the idle channel.

STEP 4

SU hashes y with the secret key k and sends the new value yk to QS to find out whether CFk of QS contains yk

STEP 5

It senses the channel that was included in the query. If the result of the sensing complies with the outcome of the Lookup operation in CFk, then SU can conclude that this channel is available

STEP 6

DB can pre-compute several cuckoo filters for each possible combination of secret keys k.

STEP 7

For spectrum opportunities the DB shares a secret key k with SU and sends the corresponding CFk to QS, also gets Channel.

2. AES:

AES (acronym of Advanced Encryption Standard) is a symmetric encryption algorithm. The algorithm was developed by two Belgian cryptographers Joan Daemen and Vincent Rijmen. AES was designed to be efficient in both hardware and software, and supports a block length of 128 bits and key lengths of 128, 192, and 256 bits

DESCRIPTION:

1. Key Expansions: round keys are derived from the cipher key using Rijndael's key schedule. AES requires a separate 128-bit round key block for each round plus one more.

2. Initial Round

a. AddRoundKey: each byte of the state is combined with a block of the round key using bitwise XOR.

3. Rounds

a. SubBytes: a non-linear substitution step where each byte is replaced with another according to a lookup table.

b. ShiftRows: a transposition step where the last three rows of the state are shifted cyclically a certain number of steps.

c. MixColumns: a mixing operation which operates on the columns of the state, combining the four bytes in each column.

d. AddRoundKey

4. Final Round (no MixColumns)

a. SubBytes

STEPS IN ADVANCED ENCRYPTION STANDARD:

STEP 1

Derive the set of round keys from the cipher key

STEP 2

Initialize the state array with the block data

STEP 3

Add the initial round key to the starting state array

STEP 4

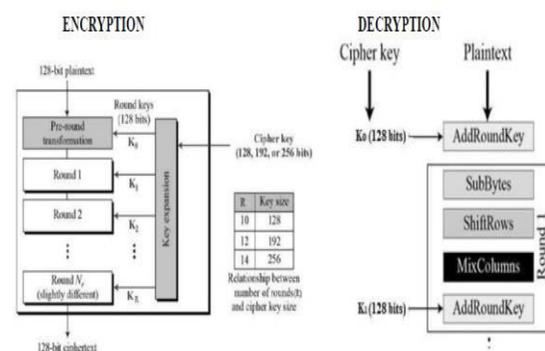
Perform nine rounds of state manipulation

STEP 5

Perform the tenth and final round of state manipulation

STEP 6

Copy the final state array out as the encrypted data



V.RESULTS

Performance Evaluation Results

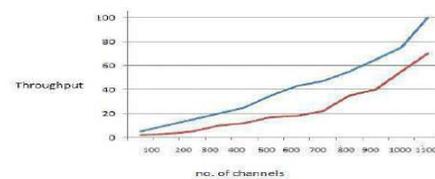


Fig. 6. Briefs about performance

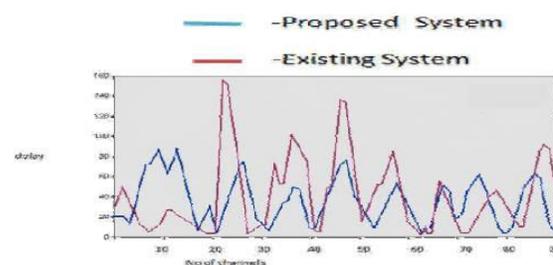


Fig. 7. Delay graph

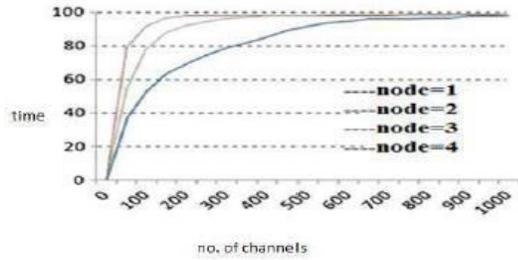


Fig.8. Node Utilisation

Fig.9. describes about registration for user in network using necessary details like name and password.

Fig.10. briefs how the server setin to network and tracks about user entering details and invader's fake details.

GET ATTACKER DETAILS

VIEW NETWORK ATTACKERS

IP	Fake Userid	Fake Password	Attack Time	Details
192.168.225.23	123	abde	2018/02/22 05:56:16	VIEW DETAILS
192.168.225.23	123	asd	2018/02/22 05:56:17	VIEW DETAILS

Fig.11. shows invader's fake details.

IP ADDRESS LOCATION TRACKING DETAILS

Network Server: Phoenix City: AS55236 Reference No: Infocentrum 1101006 Diberaparam

Country: Country Code: IN

Latitude: Longitude: 13.082601999999999 80.2797184

Gateway: Organization: 197.20.222.18 Jio

Region Name: Time Zone: Yambli Wachu Asia/Kolkata

Fig.12. shows invader's network details.

TRACK ATTACKER LOCATION

VIEW ATTACKER LOCATION ON MAP

IP	Latitude Location	Longitude Location	Address	View
192.168.225.23	13.082601999999999	80.2797184	GET ADDRESS	VIEW MAP
192.168.225.23	13.082601999999999	80.2797184	GET ADDRESS	VIEW MAP

Fig.13. shows geographical info. of invader.

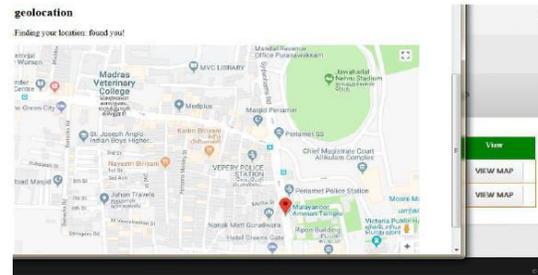


Fig.14. provides a view of exact location of invader.

BLOCK ATTACKER IP

VIEW ATTACKER LOCATION ON MAP

IP	Username Try	Password Try	Block
192.168.225.23	123	abde	BLOCKED
192.168.225.23	123	asd	BLOCKED

Fig.15. provides a status of invader blocking

TOOLS DESCRIPTION:

- Front End - HTML, J2EE
- Server side Script - Java Server Pages.
- Database - My sql

Database Connectivity - JDBC.

Java:

Java technology is both a programming language and a platform.

The Java Programming Language

The Java programming language is a high-level language that can be characterized by all of the following buzzwords:

- Simple •
- Architecture neutral •
- Object oriented •
- Portable •
- Distributed •



- High performance
- Interpreted
- Multithreaded
- Robust
- Dynamic
- Secure

Java is developed by Sun Microsystems, later acquired by Oracle Corporation, that provides a system for developing application software and deploying it in a cross-platform computing environment. Java is used in a wide variety of computing platforms from embedded devices and mobile phones to enterprise servers and supercomputers. While less common, Java applets run in secure, sand boxed environments to provide many features of native applications and can be embedded in HTML pages. Writing in the Java programming language is the primary way to produce code that will be deployed as byte code in a Java Virtual Machine (JVM); byte code compilers are also available for other languages, including Ada, JavaScript, Python, and Ruby. In addition, several languages have been designed to run natively on the JVM, including Scala, Clojure and Groovy. Java syntax borrows heavily from C and C++, but object-oriented features are modeled after Smalltalk and Objective-C.[11] Java eschews certain low-level constructs such as pointers and has a very simple memory model where every object is allocated on the heap and all variables of object types are references. Memory management is handled through integrated automatic garbage collection performed by the JVM.

Back End:

MySQL, the most popular Open Source SQL database management system, is developed, distributed, and supported by Oracle Corporation. The MySQL Web site (<http://www.mysql.com/>) provides the latest information about MySQL software.

MySQL is a database management system:

A database is a structured collection of data. It may be anything from a simple shopping list to a picture gallery or the vast amounts of information in a corporate network. To add, access, and process data stored in a computer database, we need a database management system such as MySQL Server. Since computers are very good at handling large amounts of data, database management systems play a central role in computing, as standalone utilities, or as parts of other applications.

MySQL databases are relational:

A relational database stores data in separate tables rather than putting all the data in one big storeroom. The database structures are organized into physical files optimized for speed. The logical model, with objects such as databases, tables, views, rows, and columns, offers a flexible programming environment. The SQL part of “MySQL” stands for “Structured Query Language”. SQL is the most common standardized language used to access databases. Depending on your programming environment, you might enter SQL directly (for example, to generate reports), embed SQL statements into code written in another language, or use a language-specific API that hides the SQL syntax. SQL is defined by the ANSI/ISO SQL Standard. The SQL standard

has been evolving since 1986 and several versions exist. “SQL-92” refers to the standard released in 1992, “SQL:1999” refers to the standard released in 1999, and “SQL:2003” refers to the current version of the standard. We use the phrase “the SQL standard” to mean the current version of the SQL Standard at any time.

MySQL software is Open Source:

The MySQL software uses the GPL (GNU General Public License), <http://www.fsf.org/licenses/>, to define what you may and may not do with the software in different situations. If you feel uncomfortable with the GPL or need to embed MySQL code into a commercial application, you can buy a commercially licensed version from us. See the MySQL Licensing Overview for more information (<http://www.mysql.com/company/legal/licensing/>).

MySQL Server works in client/server or embedded systems:

The MySQL Database Software is a client/server system that consists of a multi-threaded SQL server that supports different backends, several different client programs and libraries, administrative tools, and a wide range of application programming interfaces (APIs).

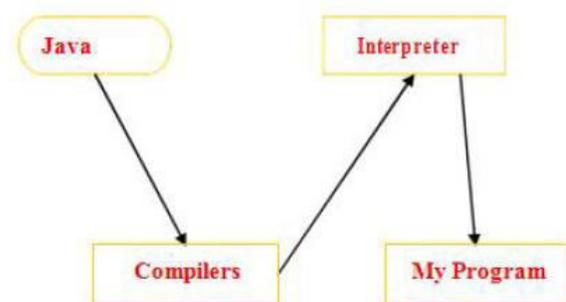
A large amount of contributed MySQL software is available:

MySQL Server has a practical set of features developed in close cooperation with our users. It is very likely that your favorite application or language supports the MySQL Database Server.

JDBC

In an effort to set an independent database standard API for Java; Sun Microsystems

developed JavaDatabase Connectivity, or JDBC. JDBC offers a generic SQL database access mechanism that provides a consistent interface to a variety of RDBMSs. This consistent interface is achieved through the use of “plug-in” database connectivity modules, or drivers. If a database vendor wishes to have JDBC support, he or she must provide the driver for each platform that the database and Java run on. To gain a wider acceptance of JDBC, Sun based JDBC’s framework on ODBC. As you discovered earlier in this chapter, ODBC has widespread support on a variety of platforms. Basing JDBC on ODBC will allow vendors to bring JDBC drivers to market much faster than developing a completely new connectivity solution. JDBC was announced in March of 1996. It was released for a 90 day public review that ended June 8, 1996. Because of user input, the final JDBC v1.0 specification was released soon after. The remainder of this section will cover enough information about JDBC for you to know what it is about and how to use it effectively. This is by no means a complete overview of JDBC. That would fill an entire book.



V. CONCLUSION

The system achieves the location and data privacy preservation of secondary user in wireless cognitive networks. It also stores the

data and geo-location of secondary user's information. Even if one of the coordinates is intentionally revealed by a SU, its location is still indistinguishable from remaining possible locations. This entity, referred to as query server (QS), has a dedicated high throughput link with DB. QS is used to guarantee computational location privacy while reducing the computational and communication overhead especially on SU's side.

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